Managing Dynamic Shared State

448430 Spring 2009 3/30/2009 Kyoung Shin Park Multimedia Engineering Dankook University

Overview

- Dynamic Shared State
- □ Consistency-Throughput Tradeoff
- Centralized/Shared Repository
- □ Frequent State Regeneration (Blind Broadcast)
- Dead Reckoning

What is Dynamic Shared State?

- The dynamic information that multiple hosts must maintain about the NVE
- Accurate dynamic shared state is fundamental to creating realistic virtual environments. It is what makes a NVE "multi-user".
- Management is one of the most difficult challenges facing the net-VE designer. The trade off is between resources and realism.

Network Latency Problem



Consistency-Throughput Tradeoff



- "It is impossible to allow dynamic shared state to change frequently and guarantee that all hosts simultaneously access identical versions of that state."
- We can have either a dynamic world or a consistent world, but not both.

Another Example

- □ Player A moves and generates a location update.
- □ To ensure consistency, player A must await acknowledgments.
- If network lag is 100 ms, acknowledgment comes no earlier than 200 ms.
- □ Therefore, player A can only update his state 5 times per second.

Design Implications

- For a highly dynamic shared state, hosts must transmit more frequent data updates.
- To guarantee consistent views of the shared state, hosts must employ reliable data delivery.
- Available network bandwidth must be split between these two constraints.

Tradeoff Spectrum

System Characteristic	Absolute Consistency	High Update Rate
View consistency	Identical at all host	Determined by data received at each host
Dynamic data support	Low: Limited by consistency protocol	High: Limited only by available bandwidth
Network infrastructure requirements	Low latency, high reliability, limited variability	Heterogeneous network possible
Number of participants supported	Low	Potentially high

Managing Shared States



Centralized / Shared Repositories

- Description: Maintain shared state data in a centralized location.
- Protect shared states via a lock manager to ensure ordered writes.
- □ Three Techniques
 - Shared File Directory
 - Repository in Server Memory
 - Distributed Repository (Virtual Repository)

Shared File Directory

- □ Absolute Consistency!
- Only one host can write data to the same file at a time. Must have locks.
- □ Slow!
- Does not support many users.

Shared File Directory



Server Memory

- Faster than Shared File Directory because each host uses does not have to open and close each file remotely.
- Don't have to have locks. Server arbitrates.
- □ Server crash is catastrophic.
- Maintaining constant connection may strain server resources.

Server Memory



Virtual Repository

- □ Tries to reduce bottleneck at server.
- Hosts communicate directly to each other following a protocol of information sharing. (can even tailor who you talk to).
- Better fault tolerance (Depending on protocol creating the "virtual files")
- Eventual Consistency.

Virtual Repository



Advantages of Centralized/ Shared Repositories

- □ Provides an easy programming model.
- **Guarantees information consistency.**
- No sense of data ownership is imposed; any host can update any piece of the shared state.

Disadvantages of Shared Repositories

- Data access and update operations require an unpredictable amount of time to complete.
- Requires considerable communications overhead due to reliable data delivery.
- □ Vulnerable to single point failure.
- **D** Push systems may send info where it is not needed.
- Limited number of users (else you overload the server or the network)
- One slow user can drag everyone down.

Frequent State Regeneration/Blind Broadcasts

- Owner of each state transmits the current value asynchronously and unreliably at regular intervals.
- Clients cache the most recent update for each piece of the shared state.
- Hopefully, frequent state update compensate for lost packets.
- No assumptions made on what information the other hosts have.
- □ Broadcast is sent "blind" to everyone.
- **u** Usually the entire entity state is sent.
- No acknowledgements
- □ No assurances of delivery
- □ No ordering of updates.

Why use?

- Can't afford overhead of centralized repository
- May not have demanding consistency requirements

Explicit Object Ownership

- With blind broadcasts, multiple hosts must not attempt to update an object at the same time.
- Each host takes explicit ownership of one piece of the shared state (usually the user's avatar).
- To update an un-owned piece of the shared state, either proxy updates or ownership transfer is employed.
- Unlike "locks" in Shared Filed servers, multiple updates are allowed until ownership is transferred
- Commonly used in online gaming (DOOM, Diablo), telesurgery, and in trying to simulate video conferencing.

Explicit Object Ownership Gaining Ownership



Explicit Object Ownership Proxy Update



Explicit Object Ownership Transferring Ownership



Reducing Broadcast Scope

- Each host sends updates to all participants in the NVE.
- Reception of extraneous updates consumes bandwidth and CPU cycles.
- Need to filter updates, perhaps at a central server (e.g. RING system) which forwards updates only to those who can "see" each other.
- VEOS epidemic approach- each host send info to specific neighbors.

Advantages of Blind Broadcasts

- Simple to implement; requires no servers, consistency protocols or a lock manager (except for filter or shared items)
- Can support a larger number of users at a higher frame rate and faster response time.

Disadvantages of Blind Broadcasts

- □ Requires large amount of bandwidth.
- Network latency impedes timely reception of updates and leads to incorrect decisions by remote hosts.
- Network jitter impedes steady reception of updates leading to 'jerky' visual behavior.
- Assumes everyone broadcasting as the same rate which may be noticeable to users if it is not the case (this may be very noticeable between local users and distant destinations).

Dead Reckoning Protocols

- Transmit state updates less frequently by using past updates to estimate the true shared state.
- □ Prediction
 - How the object's current state is computed based on previously received packets.
 - Each host estimates entity locations based on past data.
- □ Convergence
 - How the object's estimated state is corrected when another update is received.
- □ No need for central server.
- □ Sacrifices accuracy of shared state for more participants.

Dead Reckoning Illustration



Prediction Using Derivative Polynomials

- □ Zero Order (simplest)
 - x(t + ∆t) = x(t) (really state regeneration-assumes the object doesn't move)
- □ First Order (velocity)
 - $x(t + \Delta t) = x(t) + v_x \Delta t$
- □ Second Order (acceleration)
 - $x(t + \Delta t) = x(t) + v_x \Delta t + a_x (\Delta t)^2$
- □ Higher Order Approximations

Hybrid Polynomial Prediction

- Need not be absolute
- Instead of using a fixed prediction scheme, dynamically choose between first or second order based on object's history.
- Use first order when acceleration is negligible or acceleration information is unreliable (changes frequently).

Hybrid Polynomial Prediction

- Position History-Based Dead Reckoning Protocol (PHBRR) - only included position. Required the host machine to calculate velocity and acceleration based on past positions.
- Actually more accurate as "snapshot" velocities and accelerations can be misleading.

Limitations of Derivative Polynomials

- □ Why not use more terms?
 - greater bandwidth required
 - greater computational complexity
 - less accurate prediction since higher order terms are harder to estimate and errors have disparate impact
- Do not take into account capabilities or limitations of objects.

Object Specialized Prediction

- Object behavior may simplify prediction scheme.
- A plane's orientation angle is determined solely by its forward velocity and acceleration.
- □ Land based objects need only two dimensions specified.

Object Specialized Prediction

- Desired level of detail often do not need to be precise with some aspects. Do we have to accurately model the flicker of the flames of a burning vehicle or is it enough to say it is on fire.
- □ The same with smoke. Some VEs need to accurately model smoke, other do not.

Convergence Algorithms

D Tells us what to do to correct an inexact prediction:



 Trade-off between computational complexity and perceived smoothness of displayed entities

Convergence Algorithms

- □ Zero order or snap convergence
- Advantage: Simple
- Disadvantage: Poorly models real world & "Jumping" entities may distract users.



Convergence Algorithms

- Cubic Spline
- □ Advantage: Smoothest looking convergence
- Disadvantage: Computationally expensive



Convergence Algorithms

- Linear Convergence
- □ Advantage: Avoids jumping
- Disadvantage: Does not prevent "sudden "or unrealistic changes in speed or direction.



Predicted Track

Convergence Algorithms

- Choice of convergence algorithm may vary within a VE depending on the entity type and characteristics.
- □ Hybrids may be used (PHBRR)

Non-Regular Updates

- Slow update rate if prediction at remote host is within an error tolerance.
- The source host models the prediction algorithm used by the remote hosts.
- Only transmit an update when an error threshold is reached or after a timeout period (heartbeat).
- Entities must have a "heartbeat" otherwise cannot distinguish between live entities and ones that have left the system.

Advantages of Dead Reckoning

- Reduces bandwidth requirements because updates are sent less frequently.
- □ Potentially larger number of players.
- Each host does independent calculations

Disadvantages of Dead Reckoning

- □ Not all hosts share the identical state about each entity.
- Protocols are more complex to implement to develop, maintain and evaluate.
- Must customize for object behavior to achieve best results.
- □ Must have convergence to cover prediction errors.
- **Collision** detection difficult to implement.
- Poor convergence methods lead to jerky movements and distract from immersion.

Conclusions

- Shared state maintenance is governed by the Consistency
 Throughput Tradeoff.
- □ Three broad types of maintenance:
 - Centralized/Shared repository
 - Frequent State Regeneration(Blind Broadcast)
 - Dead Reckoning
- The correct choice relies on balancing many issues including bandwidth, latency, data consistency, reproducibility, and computational complexity.